

Figure 4.1 Structuring information about the family.
family(
person( tom, fox, date(7,may,1960), works(bbc,15200) ), person( ann, fox, date(9,may,1961), unemployed),
[ person( pat, fox, date(5,may,1983), unemployed), person( jim, fox, date(5,may,1983), unemployed) ] ).

(b)

(c)


Figure 4.2 Specifying objects by their structural properties: (a) any Armstrong family; (b) any family with exactly three children; (c) any family with at least three children. Structure (c) makes provision for retrieving the wife's name through the instantiation of the variables Name and Surname.


## Figure 7.10

A finite state machine (FSM) to parse an identifier.

Example 7.4 To parse the string cab3, you would make the following transitions:

| Current state: A | Input: cab3 | Scan c and go to B. |
| :--- | :--- | :--- |
| Current state: B | Input: ab3 | Scan a and go to B. |
| Current state: B | Input: b3 | Scan b and go to B. |
| Current state: B | Input: 3 | Scan 3 and go to B. |
| Current state: B | Input: | Check for final state. |

Because there is no more input and the last state is B, a final state, cab3 is a valid identifier.


## Figure 7.14

A nondeterministic FSM to parse a signed integer.

Example 7.5 You must make the following decisions to parse +203 with this nondeterministic FSM:

Current state: A Input: +203 Scan + and go to B.
Current state: B Input: 203 Scan 2 and go to B.
Current state: B Input: 03 Scan 0 and go to B.
Current state: B Input: 3 Scan 3 and go to C.
Current state: C Input: Check for final state.
Because there is no more input and you are in the final state C, you have proven that the input string +203 is a valid signed integer.


## Figure 7.17

An FSM with an empty transition to parse a signed integer.

Example 7.6 To parse 32 requires the following decisions:

| Current state: I | Input: 32 | Scan $\epsilon$ and go to F. |
| :--- | :--- | :--- |
| Current state: F | Input: 32 | Scan 3 and go to M. |
| Current state: M | Input: 2 | Scan 2 and go to M. |
| Current state: M | Input: | Check for final state. |

The transition from I to F on $\epsilon$ does not consume an input character. When you are in state I, you can do one of three things: (a) scan + and go to F, (b) scan - and go to F , or (c) scan nothing (that is, the empty string) and go to F .


Figure 4.3 An example of a non-deterministic finite automaton.


Figure 4.3 An example of a non-deterministic finite automaton.
final( s 3 ).

$$
\begin{aligned}
& \operatorname{trans}(s 1, a, s 1) \\
& \operatorname{tran} s(s 1, a, s 2) \\
& \operatorname{trans}(s 1, b, s 1) \\
& \operatorname{tran}(s 2, b, s 3) \\
& \operatorname{trans}(s 3, b, s 4) \\
& \operatorname{silent}(s 2, s 4) \\
& \operatorname{silent}(s 3, s 1)
\end{aligned}
$$



Figure 4.4 Accepting a string: (a) by reading its first symbol $X$; (b) by making a silent move.

